— MODULE working -

EXTENDS Naturals, Sequences

CONSTANTS USERIDS, SERVERS, METADATAS, IMAGES, This constant is added to allow us to assign UUIDs as blob store keys UUIDS

VARIABLES

 $databaseState, \\ blobStoreState, \\ serverStates, \end{cases}$

operations

 $vars \triangleq \langle databaseState, blobStoreState, serverStates, operations \rangle$

Strong Typing

[

Describes a database record. We need this now that it has to keep track of which image UUID it is associated with.

```
DatabaseRecord \triangleq [
metadata : MetadataVal,
imageId : UUIDVal
]
```

Describes all possible states a server can be in. Unchanged since last example. ServerStateVal $\stackrel{\Delta}{=}$

state : {
 current:
 "waiting", next: StartWrite or StartRead
 after: StartWrite
 "started_write", next: WriteBlob or FailWrite
 after: WriteBlob
 "wrote_blob", next: WriteMetadataAndReturn or FailWrite
 after: StartRead

```
"started_read", next: ReadMetadata
after: ReadMetadata, ReadMetadataAndReturnEmpty
"read_metadata" next: ReadBlobAndReturn
},
userId : UserIdVal,
metadata : MetadataVal,
imageId : UUIDVal, Need to track imageId to perform a lookup
image : ImageVal
```

This is an observability value, and we are still measuring the same thing No changes are needed $OperationValue \triangleq [type: { "READ", "WRITE" }, userId: UserIdVal, metadata : MetadataVal, image : Image Val]$

$TypeOk \triangleq$

Database state modified to hold database records $\land databaseState \in [USERIDS \rightarrow DatabaseRecord]$ Blob store uses UUIDs as keys rather than userIds $\land blobStoreState \in [UUIDS \rightarrow ImageVal]$ $\land serverStates \in [SERVERS \rightarrow ServerStateVal]$ $\land operations \in Seq(OperationValue)$

Init \triangleq

State Machine: All of the states are functions of s (server), because the only actively modeled actors in this system are our servers, but there can be multiple working simultaneously.

Writes

 $StartWrite(s) \triangleq$

 \land serverStates[s].state = "waiting" $\land \exists u \in USERIDS, m \in METADATAS, i \in IMAGES:$ \land serverStates' = [serverStates EXCEPT $![s].state = "started_write",$ Set values for the upcoming write ![s].userId = u,![s].metadata = m,![s].image = i]Record the write for observability $\land operations' = Append(operations,$ $type \mapsto$ "WRITE", $userId \mapsto u$, metadata $\mapsto m$, $image \mapsto i$ \wedge UNCHANGED $\langle databaseState, blobStoreState \rangle$ $WriteBlob(s) \stackrel{\Delta}{=}$ LET currentState $\stackrel{\Delta}{=}$ serverStates[s] IN $\land currentState.state = "started_write"$ $\wedge \exists id \in UUIDS :$ Guarantees a unique Id to simulate UUID. Note: If we run out of unset UUIDs, our system will just stop writing. We need to look out for this and ensure the set of UUIDsis large enough. \land *blobStoreState*[*id*] = "UNSET" \land blobStoreState' = [blobStoreState EXCEPT ![id] = currentState.image]Track Id to write to database \land serverStates' = [serverStates EXCEPT] $![s].state = "wrote_blob",$![s].imageId = id] \wedge unchanged $\langle databaseState, \ operations \rangle$ Writing the database is now the last part of a write operation $WriteMetadataAndReturn(s) \triangleq$ LET currentState \triangleq serverStates[s] IN $\land currentState.state = "wrote_blob"$ $\land databaseState' = [databaseState except]$![currentState.userId] = [$metadata \mapsto currentState.metadata,$ Store *imageId* in database for read $imageId \mapsto currentState.imageId$]]

```
\land serverStates' = [serverStates EXCEPT
                                     ![s].state = "waiting",
                                     ![s].userId = "UNSET".
                                     ![s].metadata = "UNSET",
                                     ![s].image = "UNSET",
                                     ![s].imageId = "UNSET"]
        \wedge UNCHANGED \langle blobStoreState, operations \rangle
FailWrite(s) \stackrel{\Delta}{=}
    \land serverStates[s].state \in { "started_write", "wrote_blob" }
    \land serverStates' = [serverStates EXCEPT
                                 ![s].state = "waiting",
                                 ![s].userId = "UNSET"
                                 ![s].metadata = "UNSET",
                                 ![s].image = "UNSET",
                                 ![s].imageId = "UNSET"]
    \wedge UNCHANGED \langle databaseState, blobStoreState, operations \rangle
```

Reads

```
StartRead(s) \stackrel{\Delta}{=}
     Reading only starts when a server is waiting
    \land serverStates[s].state = "waiting"
    \land \exists u \in USERIDS :
              serverStates' = [serverStates \ EXCEPT]
                                     ![s].state = "started_read",
                                     ![s].userId = u]
    \wedge UNCHANGED \langle databaseState, blobStoreState \rangle
    \land UNCHANGED operations
 If database record is present
ReadMetadata(s) \stackrel{\Delta}{=}
    LET currentState \stackrel{\Delta}{=} serverStates[s]
    IN
    \land currentState.state = "started_read"
     Represents reading the metadata while the database record is set
    \land databaseState[currentState.userId].metadata \neq "UNSET"
    \land serverStates' =
        [serverStates EXCEPT
            ![s].state = "read_metadata",
            ![s].metadata = databaseState[currentState.userId].metadata,
             Reads imageId from database
            ![s].imageId = databaseState[currentState.userId].imageId]
    \wedge UNCHANGED \langle databaseState, blobStoreState \rangle
```

 \wedge UNCHANGED operations

```
If database record is not present
ReadMetadataAndReturnEmpty(s) \triangleq
   LET currentState \stackrel{\Delta}{=} serverStates[s]
   IN
    \land currentState.state = "started_read"
     Represents reading the metadata while the database record is unset
    \land databaseState[currentState.userId].metadata = "UNSET"
    \land serverStates' = [serverStates EXCEPT
                                 ![s].state = "waiting"
                                 ![s].userId = "UNSET"
                                 ![s].metadata = "UNSET",
                                 ![s].image = "UNSET",
                                 ![s].imageId = "UNSET"]
    \land operations' = Append(operations,
                                Returns an empty record
                                   type \mapsto "READ",
                                   userId \mapsto currentState.userId,
                                   metadata \mapsto "UNSET",
                                   image \mapsto "UNSET"
                                ])
    \wedge UNCHANGED \langle databaseState, blobStoreState \rangle
ReadBlobAndReturn(s) \stackrel{\Delta}{=}
                           \stackrel{\Delta}{=} serverStates[s]
   Let currentState
   IN
    \land currentState.state = "read_metadata"
    \land operations' = Append(operations,
                                   type \mapsto "READ",
                                   userId \mapsto currentState.userId,
                                   metadata \mapsto currentState.metadata,
                                    Looks up image by imageId
                                   image \mapsto blobStoreState[currentState.imageId]
                               ])
    \land serverStates' = [serverStates EXCEPT
                             ![s].state = "waiting",
                             ![s].userId = "UNSET"
                             ![s].metadata = "UNSET"
                             ![s].image = "UNSET",
                             ![s].imageId = "UNSET"]
    \land UNCHANGED \langle databaseState, blobStoreState \rangle
```

Specification / Next

```
Specification / Next

Next \triangleq

For every step, pick a server and have it advance one state

\exists s \in SERVERS :

\lor StartWrite(s)

\lor WriteBlob(s) New step

\lor WriteMetadataAndReturn(s) New step

\lor FailWrite(s)

\lor StartRead(s)

\lor ReadMetadata(s) New step

\lor ReadMetadataAndReturnEmpty(s) New step

\lor ReadBlobAndReturn(s)
```

 $Spec \stackrel{\Delta}{=} Init \land \Box [Next]_{vars}$

Invariants

Note that the success criteria hasn't changed this whole time

 $ConsistentReads \triangleq$ If there are no operations, they are consistent \lor operations = $\langle \rangle$ $\forall \forall i \in 1 \dots Len(operations)$: For every read operation LET $readOp \stackrel{\Delta}{=} operations[i]$ IN \wedge readOp.type = "READ" \vee There must exist a write operation $\land \lor \exists j \in 1 \dots i :$ LET write $Op \stackrel{\Delta}{=} operations[j]$ IN \land writeOp.type = "WRITE" With the same data \land readOp.userId = writeOp.userId \land readOp.metadata = writeOp.metadata \land readOp.image = writeOp.image \lor Ignore unset reads \land readOp.metadata = "UNSET" \wedge readOp.image = "UNSET" \lor readOp.type = "WRITE" Ignore writes

This is used for model checker configuration so the simulation doesn't go on forever.

 $\begin{array}{l} StopAfter3Operations \ \triangleq \\ Len(operations) \leq 3 \end{array}$ $\begin{array}{l} StopAfter5Operations \ \triangleq \\ Len(operations) \leq 5 \end{array}$